

A New Approach to Assessing the Validity of Indels in Algorithmic Pair Alignments

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Abstract—Analysis of the structure of indels in algorithmic versus evolutionary alignments based on a set of inequalities confirms the conclusions from numerical modeling. For the more divergent sequences (PAM > 60), the tested aligning algorithm (SW) tends to increase the mean length of indels and decrease their number.

Key words: pairwise alignment, insertion, deletion

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INTRODUCTION

The sequences of symbols in biopolymers can be compared using various algorithms of pairwise alignment [1–6]. Applying any procedure, it is important to have an idea of the quality of the aligning algorithm, i.e., of how precisely the algorithmic alignment (obtained by optimization of a certain target function) can reproduce the “evolutionary alignment” of the amino acid (or nucleotide) sequences. The evolutionary alignment implies collation of the positions in the compared polymers that originate from one and the same position of their common ancestor.

Here, the “evolutionarily true” (ET) alignment is based on the data on the disposition of insertions/deletions and substitutions introduced into sequences [7, 8] in an evolutionary model [9, 10].

In a recent work [8] we used numerical modeling to assess the quality of the global version of the Smith–Waterman (SW) algorithm and compare the structure of indels and changes in SW vs. ET alignments. We found that in the algorithmic alignments the overall length of indels was markedly smaller than in the ET ones, because the number of indels decreased considerably while the mean indel length remained the same or slightly increased [8].

Here we present an alternative theoretical approach, whereby the indel characteristics for algorithmic and true alignments can be compared by their mean values without statistical processing of test sets of alignments.

THE METHOD

Basic relationships. Consider the following ratios for indel characteristics in the i -th pair of algorithmic (alg) and the evolutionarily true (et) alignments:

$v_i = n_{\text{alg}}^i / n_{\text{et}}^i$ for the number of indels,

$\lambda_i = l_{\text{alg}}^i / l_{\text{et}}^i$ for the sum length of indels,

$\alpha_i = \lambda_i / v_i$ for the mean length of an indel.

Then the corresponding averages over the set of alignments are

$$v_{\text{av}} = \frac{1}{N} \sum_{i=1 \dots N} v_i,$$

$$\lambda_{\text{av}} = \frac{1}{N} \sum_{i=1 \dots N} \lambda_i,$$

$$\alpha_{\text{av}} = \frac{1}{N} \sum_{i=1 \dots N} \lambda_i / v_i.$$

where N is the number of aligned pairs.

Now compare the sum length of indels with the product of mean indel length by the number of indels:

$$\lambda_{\text{av}} \simeq v_{\text{av}} \alpha_{\text{av}}$$

or

$$(1/N) \sum_{i=1 \dots N} \lambda_i \simeq (1/N^2) \left(\sum_{i=1 \dots N} v_i \right) \left(\sum_{i=1 \dots N} \lambda_i / v_i \right), \quad (1)$$

whereby the difference is

$$\Delta \equiv (1/N) \sum_{i=1 \dots N} \lambda_i - (1/N^2) \left(\sum_{i=1 \dots N} v_i \right) \left(\sum_{i=1 \dots N} \lambda_i / v_i \right). \quad (2)$$

It is easy to see that

Table 1

$\Delta > 0$	$\sigma_v/v_{av} > \sigma_\lambda/\lambda_{av}$	–
$\Delta > 0$	$\sigma_v/v_{av} < \sigma_\lambda/\lambda_{av}$	$\Delta v_{ij} > 0, \Delta \lambda_{ij} > 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} < 0$
$\Delta < 0$	$\sigma_v/v_{av} > \sigma_\lambda/\lambda_{av}$	$\Delta v_{ij} > 0, \Delta \lambda_{ij} > 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} < 0$ $\Delta v_{ij} > 0, \Delta \lambda_{ij} < 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} > 0$
$\Delta < 0$	$\sigma_v/v_{av} < \sigma_\lambda/\lambda_{av}$	$\Delta v_{ij} > 0, \Delta \lambda_{ij} < 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} > 0$

Table 2

Length	PAM [9]	Product αv [8]
200	60	0.88
200	100	0.84
200	200	0.77
200	300	0.77
500	60	0.92
500	100	0.89
500	200	0.82
500	300	0.82

$$(i) \left(\sum_{i=1 \dots N} v_i \right) \left(\sum_{i=1 \dots N} \alpha_i \right) =$$

$$= \sum_{i=1 \dots N} \lambda_i + \sum_{i=1 \dots N-1} \sum_{j=i+1 \dots N} (v_i \alpha_j + v_j \alpha_i),$$

$$(ii) N \sum_{i=1 \dots N} \lambda_i = \sum_{i=1 \dots N} \lambda_i + \sum_{i=1 \dots N-1} \sum_{j=i+1 \dots N} (\lambda_i + \lambda_j).$$

and (2) can be written as

$$\Delta = (1/N^2) \sum_{i=1 \dots N-1} \sum_{j=i+1 \dots N} \delta_{ij},$$

where $\delta_{ij} = (v_j \lambda_i - v_i \lambda_j)(v_i - v_j)/v_i v_j$.

Dependence of the sign of Δ on the combination of signs for differences $\Delta v_{ij} = v_i - v_j$ and $\Delta \lambda_{ij} = \lambda_i - \lambda_j$. Note that the sign of δ_{ij} is determined by the combination of inequality signs in (a) $\Delta v_{ij}/v_j \vee \Delta \lambda_{ij}/\lambda_j$, (b) $\Delta v_{ij} \vee 0$, and (c) $\Delta \lambda_{ij} \vee 0$ where \vee can be $>$ or $<$. Thus there are eight combinations of signs, whereby

$$(\Delta v_{ij}/v_j > \Delta \lambda_{ij}/\lambda_j, \Delta v_{ij} > 0, \Delta \lambda_{ij} < 0) \Rightarrow \delta_{ij} < 0, (3)$$

$$(\Delta v_{ij}/v_j < \Delta \lambda_{ij}/\lambda_j, \Delta v_{ij} < 0, \Delta \lambda_{ij} < 0) \Rightarrow \delta_{ij} < 0, (4)$$

$$(\Delta v_{ij}/v_j < \Delta \lambda_{ij}/\lambda_j, \Delta v_{ij} < 0, \Delta \lambda_{ij} > 0) \Leftrightarrow \delta_{ij} > 0. (5)$$

Since

$$\delta_{ij} > 0, \quad i, j = 1, \dots, N \Rightarrow \Delta > 0,$$

δ_{ij} in (3)–(5) can be replaced with Δ (also replacing \Leftrightarrow with \Rightarrow).

Dependence of the relation between σ_v/v_{av} and $\sigma_\lambda/\lambda_{av}$ on the relation between $(\Delta v_{ij})^2$ and $(\Delta \lambda_{ij})^2$. Since $v_{av}, \lambda_{av} > 0$, then $\sigma_v/v_{av} > \sigma_\lambda/\lambda_{av} \sim \sigma_v \lambda_{av} > \sigma_\lambda v_{av}$.

For rmsd we have

$$\sigma_a = (1/N) \left(\sum_{i=1 \dots N-1} \sum_{j=i+1 \dots N} \Delta a_{ij}^2 \right)^{1/2}.$$

Then

$$\sigma_a c_{av} =$$

$$= (1/N)^2 \left[\left(\sum_{i=1 \dots N-1} \sum_{j=i+1 \dots N} \Delta a_{ij}^2 \right) \left(\sum_{i=1 \dots N} b_i \right)^2 \right]^{1/2},$$

where $a = v, \lambda, b = \lambda, v$.

Since

$$\left(\sum_{i=1 \dots N} b_i \right)^2 = \sum_{i=1 \dots N} b_i^2 + 2 \sum_{i=1 \dots N-1} \sum_{j=i+1 \dots N} b_i b_j$$

we can state that

$$\forall i, j, k, l = 1, \dots, N: (\Delta v_{ij})^2 \lambda_k \lambda_l > (\Delta \lambda_{ij})^2 v_k v_l \Rightarrow \sigma_v \lambda_{av} > \sigma_\lambda v_{av} \quad (6)$$

or:

if for all $i, j, k, l = 1, \dots, N$ it is true that

$$(\Delta v_{ij})^2 \lambda_k \lambda_l > (\Delta \lambda_{ij})^2 v_k v_l,$$

then $\sigma_v \lambda_{av} > \sigma_\lambda v_{av}$.

Assuming that

$$((\forall i, j, k, l = 1, \dots, N: (\Delta v_{ij})^2 \lambda_k \lambda_l < (\Delta \lambda_{ij})^2 v_k v_l) \Leftrightarrow \forall i, j, k, l = 1, \dots, N: |\Delta v_{ij}| \lambda_j < |\Delta \lambda_{ij}| v_j,$$

from (3)–(6) we get

$$(\Delta v_{ij}/v_j > \Delta \lambda_{ij}/\lambda_j, \Delta v_{ij} < 0, \Delta \lambda_{ij} < 0) \Rightarrow \{\Delta > 0, \sigma_v/v_{av} < \sigma_\lambda/\lambda_{av}\}, \quad (7)$$

$$(\Delta v_{ij}/v_j > \Delta \lambda_{ij}/\lambda_j, \Delta v_{ij} < 0, \Delta \lambda_{ij} > 0) \Rightarrow \{\Delta < 0, \sigma_v/v_{av} > \sigma_\lambda/\lambda_{av}\}, \quad (8)$$

$$(\Delta v_{ij}/v_j < \Delta \lambda_{ij}/\lambda_j, \Delta v_{ij} < 0, \Delta \lambda_{ij} < 0) \Rightarrow \{\Delta < 0, \sigma_v/v_{av} \vee \sigma_\lambda/\lambda_{av}\} \quad (9)$$

which can be presented in tabular form (Table 1).

Table 3

Length	PAM	σ_v/ν_{av}	σ_v/λ_{av}	Δ	Combination
200	60	0.1929	0.2184	-0.0031	$\Delta v_{ij} > 0, \Delta \lambda_{ij} < 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} > 0$
200	100	0.2061	0.2384	-0.0057	$\Delta v_{ij} > 0, \Delta \lambda_{ij} < 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} > 0$
200	200	0.2628	0.3033	-0.0124	$\Delta v_{ij} > 0, \Delta \lambda_{ij} < 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} > 0$
200	300	0.3311	0.3619	-0.0255	$\Delta v_{ij} > 0, \Delta \lambda_{ij} < 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} > 0$
500	60	0.1211	0.1449	0.0046	$\Delta v_{ij} > 0, \Delta \lambda_{ij} > 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} < 0$
500	100	0.1355	0.162	-0.0006	$\Delta v_{ij} > 0, \Delta \lambda_{ij} < 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} > 0$
500	200	0.1895	0.2036	-0.0284	$\Delta v_{ij} > 0, \Delta \lambda_{ij} < 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} > 0$
500	300	0.2547	0.238	-0.0118	$\Delta v_{ij} > 0, \Delta \lambda_{ij} < 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} > 0$ $\Delta v_{ij} > 0, \Delta \lambda_{ij} > 0$ $\Delta v_{ij} < 0, \Delta \lambda_{ij} < 0$

RESULTS AND DISCUSSION

Using the averaged v , λ , α values from our recent work [8], we can obtain the required (mean length by number) product (Table 2) and analyze the relative characteristics of indels for various evolutionary distances as regards the trends in sign combinations (Table 3).

Non-coincidence of the Δv_{ij} and $\Delta \lambda_{ij}$ signs (right-most column) means that a higher (lower) ratio of the indel numbers goes together with a lower (higher) ratio of the mean indel lengths (i.e., indel lengths are not the same in the algorithmic and the ET alignments). Judging by the increase in the Δ absolute value, this tendency is pronounced in all cases but two: at $L = 500$ and PAM = 60 the signs agree and the mean indel lengths may be equal; at $L = 500$ and PAM = 300 the assessment is inconclusive. In any case, this analysis shows that for the more divergent sequences (less than 55% identity) the increase of the sum length of indels in the algorithmic alignment is caused by overestimation of the mean indel length and underestimation of the indel number as compared with the true alignment. That is, the SW algorithm of pairwise alignment tends to alter the structure of indels.

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